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Recitation 1: Background on Probability Theory
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## 1 Basic theory

### 1.1 Definitions and background

Probability is always defined over a set of possible outcomes $\Omega$ of a (discrete) random experiment. For example, the possible outcomes for a single coin flip are heads $(H)$ or tails $(T)$, and so we have $\Omega=\{H, T\}$. We call $\Omega$ the sample space of the experiment. On the other hand, if our random experiment consists of flipping a coin $n$ times independently, then

$$
\Omega=\left\{\left(\omega_{1}, \ldots, \omega_{n}\right): \omega_{i} \in\{H, T\}\right\}
$$

A probability distribution over $\Omega$ is a function $p: \Omega \rightarrow \mathbb{R}_{\geq 0}$ such that $\sum_{x \in \Omega} p(x)=1$. That is, the sum of all $p_{i}$ over all possible outcomes in $\Omega$ is 1 .

An event is any subset $A \subseteq \Omega$. The probability of an event $A$ is

$$
\underset{p}{\operatorname{Pr}}[A]=\sum_{x \in A} p(x) .
$$

In words, we can define the probability of an event in a uniform distribution as

$$
\operatorname{Pr}[\text { event happens }]=\frac{\text { number of ways it can happen }}{\text { total number of outcomes }}
$$

Note: We will often just write $\operatorname{Pr}$ instead of $\operatorname{Pr}$ when the distribution $p$ is clear from context.

Two events $A, B \subseteq \Omega$ are said to be independent, if $\operatorname{Pr}[A \cap B]=\operatorname{Pr}[A] \cdot \operatorname{Pr}[B]$.
Let look at an example involving coin tosses.

## Example 1: Tossing a coin

Recall the experiment of tossing a coin $n$ times: $\Omega=\left\{\left(\omega_{1}, \ldots, \omega_{n}\right): \omega_{i} \in\{H, T\}\right\}$. The event that the first flip is heads is represented as the set

$$
A_{1, H}=\left\{\left(H, \omega_{2}, \ldots, \omega_{n}\right): \omega_{i} \in\{H, T\}\right\}
$$

and similarly the event that the first flip is tails is

$$
A_{1, T}=\left\{\left(T, \omega_{2}, \ldots, \omega_{n}\right): \omega_{i} \in\{H, T\}\right\}
$$

## Example 2: Tossing a coin (continued...)

We can define the events $A_{i, H}$ for the $i$-th flip to be heads, and $A_{i, T}$ for tails. If we assume that the coin is fair ${ }^{1}$ and each toss is independent, then we have that for fixed $\omega_{1}, \cdots \omega_{n}$ :

$$
\begin{aligned}
p\left(\left(\omega_{1}, \ldots, \omega_{n}\right)\right) & =\operatorname{Pr}\left[A_{1, \omega_{1}} \cap \ldots \cap A_{n, \omega_{n}}\right] \\
& =\operatorname{Pr}\left[A_{1, \omega_{1}}\right] \ldots \operatorname{Pr}\left[A_{n, \omega_{n}}\right] / / \text { by independence } \\
& =\frac{1}{2} \cdot \ldots \cdot \frac{1}{2} / / \text { by fairness } \\
& =\frac{1}{2^{n}}
\end{aligned}
$$

[^0]A (real-valued) random variable is a function $X: \Omega \rightarrow \mathbb{R}$. In Example 1 , the number of heads is a random variable represented by the function

$$
X\left(\left(\omega_{1}, \ldots, \omega_{n}\right)\right)=\sum_{i=1}^{n} \omega_{i}
$$

Two discrete real-valued random variables $X, Y$ are said to be independent if

$$
\operatorname{Pr}[X=x, Y=y]=\operatorname{Pr}[X=x] \cdot \operatorname{Pr}[Y=y]
$$

for any $x, y \in \mathbb{R}$.
The random variables $X_{1}, \ldots, X_{n}$ are said to be (jointly) independent if

$$
\operatorname{Pr}\left[X_{1}=x_{1}, \ldots, X_{n}=x_{n}\right]=\prod_{i=1}^{n} \operatorname{Pr}\left[X_{i}=x_{i}\right]
$$

for any $x_{1}, \ldots, x_{n}$.
Note: $X_{1}, \ldots, X_{n}$ can be pairwise independent without being jointly independent!

## Example 3: Jointly independent random variables

Let $X_{i}$ be the random variable that is 1 if the $i$-th fair coin landed heads and 0 otherwise. That is,

$$
X_{i}=\left\{\begin{array}{l}
1, \text { if } i \text { th coin lands } H \\
0, \text { if } i \text { th coin lands } T
\end{array}\right.
$$

The random variables $X_{1}, \ldots, X_{n}$ are jointly independent.

## Example 4: Pairwise independent but not jointly independent random variables

Let $\Omega=\left\{\left(\omega_{1}, \omega_{2}\right): \omega_{i} \in\{0,1\}\right\}$. Consider the uniform distribution $p$ over $\Omega$. Let $X\left(\omega_{1}, \omega_{2}\right)=\omega_{1}$ be the outcome of the first flip, $Y\left(\omega_{1}, \omega_{2}\right)=\omega_{2}$ be the outcome of the second flip, and $Z\left(\omega_{1}, \omega_{2}\right)=\omega_{1} \oplus \omega_{2}$ be the XOR of both flips. For all $x, y, z \in\{0,1\}$, we have

$$
\begin{gathered}
\operatorname{Pr}[X=x]=\operatorname{Pr}[Y=y]=\operatorname{Pr}[Z=z]=\frac{1}{2}, \text { and } \\
\operatorname{Pr}[X=x, Y=y]=\operatorname{Pr}[X=x, Z=z]=\operatorname{Pr}[Y=y, Z=z]=\frac{1}{4},
\end{gathered}
$$

which implies pairwise independence. However, we have

$$
\operatorname{Pr}[X=0, Y=0, Z=1]=0 \neq \frac{1}{8}=\operatorname{Pr}[X=0] \cdot \operatorname{Pr}[Y=0] \cdot \operatorname{Pr}[Z=1]
$$

showing that these random variables are not jointly independent.

### 1.2 Law of total probability.

The law of total probability states that if we have events $A_{1}, A_{2}, \ldots, A_{n}$ which partition the sample space (i.e., $\Omega$ is a disjoint union of these events), and $B$ is any event, then

$$
\operatorname{Pr}[B]=\sum_{i=1}^{n} \operatorname{Pr}\left[B \cap A_{i}\right]
$$

Note: the law of total probability is also valid if we have a countably infinite partition into events $A_{1}, A_{2}, \ldots, A_{n}, \ldots$, in which case

$$
\operatorname{Pr}[B]=\sum_{i=1}^{\infty} \operatorname{Pr}\left[B \cap A_{i}\right]
$$

### 1.3 Conditional probability.

Conditioning on an event means assuming examining what happens given another event occurs. For example, we can condition on the probability of it raining tomorrow given that it's raining today.
Formally, the probability of event $A$ conditioned on event $B$ is defined as

$$
\operatorname{Pr}[A \mid B]=\frac{\operatorname{Pr}[A \cap B]}{\operatorname{Pr}[B]}
$$

In words, this formula captures the probability that both events happen, subject to $B$ happening. The intuition is that we focus only on the part of our sample space $\Omega$ in which $B$ happens (i.e. the subset $B \subset \Omega)$. We can then think of $B$ as our new sample space. Then, the part of $A$ that matters is only the intersection $A \cap B$. However, if we make $B$ our new sample space, then we need to normalize by the probability that $B$ happens (in $\Omega$ ), because $B$ doesn't necessarily have the full probability mass within $\Omega$. This is where the denominator $\operatorname{Pr}[B]$ comes from in the formula.

To see that this is indeed a valid probability, note that

$$
\operatorname{Pr}[A \mid B]+\operatorname{Pr}[\bar{A} \mid B]=\frac{\operatorname{Pr}[A \cap B]+\operatorname{Pr}[\bar{A} \cap B]}{\operatorname{Pr}[B]}=1
$$

From the above, we get Bayes' rule, one of the most important formulas in probability:

$$
\operatorname{Pr}[A \mid B]=\frac{\operatorname{Pr}[B \mid A] \operatorname{Pr}[A]}{\operatorname{Pr}[B]}
$$

Additionally, using conditional probability, we can obtain a formula for the probability of an intersection of events, even if the events are not independent!

$$
\operatorname{Pr}\left[A_{1} \cap A_{2} \cdots \cap A_{n}\right]=\operatorname{Pr}\left[A_{1}\right] \operatorname{Pr}\left[A_{2} \mid A_{1}\right] \ldots \operatorname{Pr}\left[A_{n} \mid A_{1} \cap A_{2} \cap \cdots \cap A_{n-1}\right]
$$

The above is often called the "chain rule" of conditional probability.

### 1.4 Union bound.

For events $A_{1}, A_{2} \subseteq \Omega$, we have

$$
\operatorname{Pr}\left[A_{1} \cup A_{2}\right]=\operatorname{Pr}\left[A_{1}\right]+\operatorname{Pr}\left[A_{2}\right]-\operatorname{Pr}\left[A_{1} \cap A_{2}\right] .
$$

This is a special case of the inclusion-exlusion principle. In particular,

$$
\operatorname{Pr}\left[A_{1} \cup A_{2}\right] \leq \operatorname{Pr}\left[A_{1}\right]+\operatorname{Pr}\left[A_{2}\right],
$$

and by induction, for $A_{1}, \cdots, A_{n} \subseteq \Omega$,

$$
\operatorname{Pr}\left[\bigcup_{i=1}^{n} A_{i}\right] \leq \sum_{i=1}^{n} \operatorname{Pr}\left[A_{i}\right] .
$$

This is called the union bound. This technique is commonly used when we want to provide a bound on the probability that at least one event happens, from a family of events. We upper bound this probability by the sum of the probabilities of the individual events. As implied by the inclusion-exclusion principle, this is tight when all the $A_{i}$ are disjoint.

### 1.5 Expectation.

For a discrete real-valued random variable $X$ taking possible values $x_{1}, \ldots, x_{n} \in \mathbb{R}$, the expectation is defined as

$$
\mathbb{E}[X]=\sum_{i=1}^{n} \operatorname{Pr}\left[X=x_{i}\right] \cdot x_{i}
$$

You can think of the expectation as being the "average" value that $X$ will take on.

## Example 5: Expectation of a coin toss

Let $X$ be the random variable that is 1 if a fair coin landed heads and 0 otherwise.

$$
\mathbb{E}[X]=\left(\frac{1}{2}\right) \cdot 1+\left(\frac{1}{2}\right) \cdot 0=\frac{1}{2} .
$$

## Linearity of Expectation

Given random variables $X_{1}, \ldots, X_{n}$ and $X=\sum_{i=1}^{n} X_{i}$, we have

$$
\mathbb{E}[X]=\mathbb{E}\left[\sum_{i=1}^{n} X_{i}\right]=\sum_{i=1}^{n} \mathbb{E}\left[X_{i}\right]
$$

In words, the expected value of the sum of random variables is equal to the sum of the expected values. A very important takeaway from this result is that it holds even if the random variables are not independent. This will be used frequently when we have to find the expected value of a sum of random variables when they might not be independent.

## Example 6: Linearity of Expectation for a coin toss

Let $X_{i}$ be the random variable that is 1 if the $i$-th fair coin landed heads and 0 otherwise.
Let $X=\sum_{i=1}^{n} X_{i}$.

$$
\mathbb{E}[X]=\mathbb{E}\left[\sum_{i=1}^{n} X_{i}\right]=\sum_{i=1}^{n} \mathbb{E}\left[X_{i}\right]=\frac{n}{2}
$$

Which matches our intuition that, on average, the number of times that the coin comes up heads is roughly half the number of trials.

## Multiplicativity of expectation under independence

Another cool property of expectation is that the expectation of a product of independent variables is the product of individual expectations:

$$
\mathbb{E}[X Y]=\mathbb{E}[X] \mathbb{E}[Y] .
$$

To see this, it is easiest to start manipulating the right side. Suppose $X$ can take values in $S$ and $Y$ can take values in $T$, and let $W=\{x y: x \in S, y \in T\}$. Then we have

$$
\begin{aligned}
\mathbb{E}[X] \mathbb{E}[Y] & =\sum_{x \in S} \sum_{y \in T} \operatorname{Pr}[X=x] \cdot \operatorname{Pr}[Y=y] \cdot x \cdot y \\
& =\sum_{x \in S} \sum_{y \in T} \operatorname{Pr}[X=x, Y=y] \cdot x \cdot y \\
& =\sum_{a \in W} \sum_{(x, y) \in S \times T: x y=a} \operatorname{Pr}[X=x, Y=y] \cdot a \\
& =\sum_{a \in W} \operatorname{Pr}[X Y=a] \cdot a=\mathbb{E}[X Y] .
\end{aligned}
$$

### 1.6 Variance.

For a discrete real-valued random variable $X$, the variance is defined as

$$
\operatorname{Var}[X]=\mathbb{E}\left[(X-\mathbb{E}[X])^{2}\right]
$$

Intuitively, the variance captures how far the random variable is from its expectation in a squared, expected sense. Note that this can be alternatively expressed as

$$
\begin{aligned}
\mathbb{E}\left[(X-\mathbb{E}[X])^{2}\right] & =\mathbb{E}\left[X^{2}-2 X \mathbb{E}[X]+\mathbb{E}[X]^{2}\right] \\
& =\mathbb{E}\left[X^{2}\right]-2 \mathbb{E}[X \mathbb{E}[X]]+\mathbb{E}[X]^{2} \\
& =\mathbb{E}\left[X^{2}\right]-2 \mathbb{E}[X]^{2}+\mathbb{E}[X]^{2} \\
& =\mathbb{E}\left[X^{2}\right]-\mathbb{E}[X]^{2} .
\end{aligned}
$$

## Linearity of variance under pairwise independence.

An important property of the variance is that it is additive when the summands are pairwise independent random variables. That is, if $X_{1}, \ldots, X_{n}$ are pairwise independent random variables, we have

$$
\operatorname{Var}\left[\sum_{i=1}^{n} X_{i}\right]=\sum_{i=1}^{n} \operatorname{Var}\left[X_{i}\right]
$$

To see this, note that

$$
\begin{aligned}
\operatorname{Var}\left[\sum_{i=1}^{n} X_{i}\right] & =\mathbb{E}\left[\left(\sum_{i=1}^{n} X_{i}\right)^{2}\right]-\left(\sum_{i=1}^{n} \mathbb{E}\left[X_{i}\right]\right)^{2} \\
& =\sum_{i=1}^{n} \mathbb{E}\left[X_{i}^{2}\right]+2 \sum_{i<j} \mathbb{E}\left[X_{i} X_{j}\right]-\sum_{i=1}^{n} \mathbb{E}\left[X_{i}\right]^{2}-2 \sum_{i<j} \mathbb{E}\left[X_{i}\right] \mathbb{E}\left[X_{j}\right] \\
& =\sum_{i=1}^{n} \mathbb{E}\left[X_{i}^{2}\right]-\sum_{i=1}^{n} \mathbb{E}\left[X_{i}\right]^{2} \\
& =\sum_{i=1}^{n} \operatorname{Var}\left[X_{i}\right]
\end{aligned}
$$

where we used the fact that $\mathbb{E}[X Y]=\mathbb{E}[X] \mathbb{E}[Y]$ for independent $X, Y$.

## 2 Concentration inequalities

Concentration inequalities are tools that allow us to bound the probability with which a random variable can be far from its expectation. There is a vast number of concentration inequalities corresponding to the different assumptions on the random variable.
For example, if a random variable is a sum of many independent random variables, intuitively it seems very (exponentially in the number of summands) unlikely for all individual random variables in the sum to conspire to bring the value of the sum away from its expectation. As we'll see below, in such a situation we in fact have theorems saying that deviating from the expectation is exponentially unlikely, as one intuitively expects.

### 2.1 Markov's inequality.

Let $Y$ be a discrete random variable taking non-negative values in the set $S$. Then for any $a>0$,

$$
\operatorname{Pr}[Y \geq a] \leq \frac{\mathbb{E}[Y]}{a}
$$

A nice feature of this inequality is that it only depends on the expectation of the random variable. Proof.

$$
\begin{aligned}
& \mathbb{E}[Y]=\sum_{y \in S} y \cdot \operatorname{Pr}[Y=y]=\sum_{y \in S, y<a} y \cdot \operatorname{Pr}[Y=y]+\sum_{y \in S, y \geq a} y \cdot \operatorname{Pr}[Y=y] \\
& \geq \sum_{y \in S, y \geq a} y \cdot \operatorname{Pr}[Y=y] \geq \sum_{y \in S, y \geq a} a \cdot \operatorname{Pr}[Y=y]=a \cdot \operatorname{Pr}[Y \geq a]
\end{aligned}
$$

This is tight when $Y$ is $a$ with probability 1 . Markov's inequality is important because it ties the probability of a random variable being greater than some threshold to the expected value of the random variable. What's not obvious though is that it can also be extended to prove much more powerful inequalities.

### 2.2 Chebyshev's inequality.

Let $X$ be a random variable with expected value $\mu$ and strictly positive variance $\sigma^{2}$. Then for all real $k>0$ :

$$
\operatorname{Pr}[|X-\mu| \geq k] \leq \frac{\sigma^{2}}{k^{2}}
$$

What this is saying is that the probability that $X$ is a distance from the mean is related directly to the variance and inversely to the squared distance. In general, Chebyshev's inequality provides us with a stronger bound than Markov's inequality because we utilize the variance of the random variable.
Proof. Since $(X-\mu)^{2}$ is a nonnegative random variable, by Markov's inequality we get

$$
\begin{aligned}
& \operatorname{Pr}\left[(X-\mu)^{2} \geq k^{2}\right] \leq \frac{\mathbb{E}\left[(X-\mu)^{2}\right]}{k^{2}} \\
& \quad \operatorname{Pr}[|X-\mu| \geq k] \leq \frac{\sigma^{2}}{k^{2}}
\end{aligned}
$$

### 2.3 Chernoff bounds.

Suppose $X_{1}, \ldots, X_{n}$ are independent random variables taking values in $\{0,1\}$. Let $X$ denote their sum and let $\mu=\mathbb{E}[X]$ denote the sum's expected value. Then for any $\beta>0$,

- $\operatorname{Pr}[X>(1+\beta) \mu]<e^{-\beta^{2} \mu / 3}$, for $0<\beta<1$
- $\operatorname{Pr}[X>(1+\beta) \mu]<e^{-\beta \mu / 3}$, for $\beta>1$
- $\operatorname{Pr}[X<(1-\beta) \mu]<e^{-\beta^{2} \mu / 2}$, for $0<\beta<1$

This allows us to get an even tighter bound because we can use the fact that the random variables exhibit full mutual independence. Note that this is a stronger assumption than pairwise independence! There are groups of random variables which are all pairwise independent but which are not mutually independent.

## 3 Example problems

### 3.1 Problems pertaining to random variables.

1. For each of the following distributions, compute their expectation and variance:
(1) Uniform in $\{1,2, \ldots, n\}$,
(2) Bernoulli ${ }^{1}$ with success probability $p$.

## Solution.

(1) We have

$$
\mathbb{E}[X]=\sum_{i=1}^{n} \frac{1}{n} \cdot i=\frac{1}{n} \cdot \frac{n(n+1)}{2}=\frac{n+1}{2}
$$

and

$$
\begin{aligned}
\operatorname{Var}[X] & =\mathbb{E}\left[X^{2}\right]-(\mathbb{E}[X])^{2}=\sum_{i=1}^{n} \frac{1}{n} \cdot i^{2}-\left(\frac{n+1}{2}\right)^{2} \\
& =\frac{1}{n} \cdot \frac{n(n+1)(2 n+1)}{6}-\frac{(n+1)^{2}}{4}=\frac{n^{2}-1}{12} .
\end{aligned}
$$

$$
\begin{equation*}
\mathbb{E}[X]=p \cdot 1+(1-p) \cdot 0=p \tag{2}
\end{equation*}
$$

and

$$
\operatorname{Var}[X]=\mathbb{E}\left[X^{2}\right]-(\mathbb{E}[X])^{2}=p \cdot 1^{2}+(1-p) \cdot 0^{2}-p^{2}=p(1-p)
$$

2. Suppose Alice flips 6 fair coins. What is the probability that result is three heads and three tails? Suppose furthermore that Alice has to pay $\$ 1$ to flip 6 coins. What is the expected number of dollars she must pay until she sees the result of three heads and three tails?
Solution. The probability space can be represented as $\Omega=\left\{\left(a_{1}, \ldots, a_{6}\right): a_{i} \in\{0,1\}\right\}$. The event of getting three heads is then $A=\left\{\left(a_{1}, \ldots, a_{6}\right): \sum_{i=1}^{6} a_{i}=3\right\}$. Since every possible choice $\left(a_{1}, \ldots, a_{6}\right)$ has the same probability $\frac{1}{2^{6}}$, we have

$$
\operatorname{Pr}[A]=\frac{|A|}{2^{6}}=\frac{\binom{6}{3}}{2^{6}}=\frac{5}{16}
$$

For the second part, we're in the following general situation: we have a Bernoulli (i.e., $\{0,1\}$ ) random variable $X$ such that $\operatorname{Pr}[X=1]=p$, and we sample independent copies $X_{1}, X_{2}, \ldots$ of $X$. We want to know what is the expected time $\mathbb{E}[T]$ such that $X_{T}=1$ for the first time. Well, we have

$$
\operatorname{Pr}[T=t]=\operatorname{Pr}\left[X_{1}=0, \ldots, X_{t-1}=0, X_{t}=1\right]=(1-p)^{t-1} p
$$

[^1]and hence
\[

$$
\begin{aligned}
\mathbb{E}[T] & =\sum_{t=1}^{\infty} \operatorname{Pr}[T=t] t=\sum_{t=1}^{\infty}(1-p)^{t-1} p t \\
& =p \sum_{t=1}^{\infty}(1-p)^{t-1} t \\
& =p\left(\sum_{t=1}^{\infty}(1-p)^{t-1}+\sum_{t=2}^{\infty}(1-p)^{t-1}+\ldots\right) \\
& =p\left(\frac{1}{p}+(1-p) \frac{1}{p}+(1-p)^{2} \frac{1}{p}+\ldots\right) \\
& =1+(1-p)+(1-p)^{2}+\ldots=\frac{1}{p}
\end{aligned}
$$
\]

So, we get a very neat result: the expected number of independent trials until a Bernoulli random variable with probability of being 1 equal to $p$ is 1 is $\frac{1}{p}$.
Applying this to our case, the expected number of dollars will be $\frac{16}{5}$.
This calculation can be simplified using the following identity which holds whenever $T$ ranges over the natural numbers:

$$
\mathbb{E}[T]=\sum_{t=0}^{\infty} \operatorname{Pr}[T>t]
$$

3. Alice flips a fair coin $n$ times, and so does Bob. Show that the probability that they get the same number of heads is $\binom{2 n}{n} / 4^{n}$. Use your argument to verify the identity

$$
\sum_{k=0}^{n}\binom{n}{k}^{2}=\binom{2 n}{n}
$$

Solution. Let our probability space be $\Omega=\left\{\left(a_{1}, \ldots, a_{n}, b_{1}, \ldots, b_{n}\right): a_{i} \in\{0,1\}, b_{i} \in\{0,1\}\right\}$, where $a_{i}=1$ if the $i$-th flip of Alice was heads and 0 otherwise, and $b_{i}=1$ if the $i$-th flip of Bob was tails, and 0 otherwise. Note that we encode heads and tails in opposite ways for Alice and Bob.
Then note that the event that they flipped the same number of heads is

$$
\begin{aligned}
A & =\left\{\left(a_{1}, \ldots, a_{n}, b_{1}, \ldots, b_{n}\right): \sum_{i=1}^{n} a_{i}=\sum_{i=1}^{n}\left(1-b_{i}\right)\right\} \\
& =\left\{\left(a_{1}, \ldots, a_{n}, b_{1}, \ldots, b_{n}\right): \sum_{i=1}^{n} a_{i}+\sum_{i=1}^{n} b_{i}=n\right\}
\end{aligned}
$$

which immediately tells us that $\operatorname{Pr}[A]=\frac{\left(2^{2}\right)}{2^{2 n}}$ as wanted.
Now, note that we could have computed the same probability with a different probability space: namely, the one where we encode heads and tails in the same way. Here $\Omega=\left\{\left(a_{1}, \ldots, a_{n}, b_{1}, \ldots, b_{n}\right): a_{i} \in\right.$ $\left.\{0,1\}, b_{i} \in\{0,1\}\right\}$, where $a_{i}=1$ if the $i$-th flip of Alice was heads and 0 otherwise, and $b_{i}=1$ if the $i$-th flip of Bob was heads, and 0 otherwise. Now we have

$$
A=\left\{\left(a_{1}, \ldots, a_{n}, b_{1}, \ldots, b_{n}\right): \sum_{i=1}^{n} a_{i}=\sum_{i=1}^{n} b_{i}\right\}
$$

We can calculate the probability by considering all the different possible numbers of heads that the two players can have (we're using the law of total probability here):

$$
\begin{aligned}
\operatorname{Pr}[A] & =\sum_{k=0}^{n} \operatorname{Pr}\left[A \cap \sum_{i=1}^{n} a_{i}=k\right] \\
& =\sum_{k=0}^{n} \operatorname{Pr}\left[\sum_{i=1}^{n} a_{i}=\sum_{i=1}^{n} b_{i}=k\right] \\
& =\sum_{k=0}^{n} \operatorname{Pr}\left[\sum_{i=1}^{n} a_{i}=k\right] \operatorname{Pr}\left[\sum_{i=1}^{n} b_{i}=k\right] \\
& =\sum_{k=0}^{n} \frac{\binom{n}{k}}{2^{n}} \frac{\binom{n}{k}}{2^{n}} \\
& =\frac{\sum_{k=0}^{n}\binom{n}{k}}{4^{n}}
\end{aligned}
$$

Comparing the two expressions, we get the desired identity.

### 3.2 Problems pertaining to concentration bounds

1. Let's say that we flip a biased coin that lands heads with probability $\frac{1}{3}$ a total of $n$ times. Use Chernoff bounds to determine a value of $n$ such that the probability of getting more than half of the flips heads is less than $\frac{1}{1000}$.
Solution. Let $X_{i}$ be a random variable that is 1 if the $i$-th flip landed heads and 0 otherwise. If we denote $X=\sum_{i=1}^{n} X_{i}$, we want to find the smallest $n$ such that $\operatorname{Pr}\left[X>\frac{n}{2}\right]<\frac{1}{1000}$.
Note that $\mu=\mathbb{E}[X]=\sum_{i=1}^{n} \mathbb{E}\left[X_{i}\right]=\sum_{i=1}^{n} \frac{1}{3}=\frac{n}{3}$. Applying Chernoff bounds from the previous section with $\beta=\frac{1}{2}$ we get

$$
\begin{aligned}
& \operatorname{Pr}\left[X>\frac{3}{2} \mu\right]<e^{-(1 / 2)^{2} \mu / 3} \\
\Leftrightarrow & \operatorname{Pr}\left[X>\frac{n}{2}\right]<e^{-n / 36}
\end{aligned}
$$

So for $e^{-n / 36}<1 / 1000 \Leftrightarrow n>36 \log 1000 \approx 250$ we have the required bound.
2. Bar the bear decides he wants to manage beehives in his old age. He's just received $k$ bees that he wants to allocate to his $n$ beehives. Since Bar is old, he often loses count when trying to allocate the bees to beehives. He decides to just allocate the bees randomly to his hives. That is, for each bee, he chooses a beehive uniformly at random. Help Bar prove that his strategy yields an approximately uniform distribution of bees with high probability.
(a) Let $X_{i}$ be the number of bees in the $i$-th beehive. Compute $E\left[X_{i}\right]$.

Solution. Let $Y_{j i}$ be 1 if the $j$-th bee is allocated to the $i$-th beehive, and 0 otherwise. We have $E\left[Y_{j i}\right]=\operatorname{Pr}[j$-th bee is put into $i$-th beehive $]=1 / n$. Then $X_{i}=\sum_{j=1}^{k} Y_{j i}$, so $E\left[X_{i}\right]=\sum_{j=1}^{k} E\left[Y_{j i}\right]=\sum_{j=1}^{k} 1 / n=k / n$.
(b) Show that $X_{i}$ and $X_{j}$ are not independent.

Solution. We see that $\operatorname{Pr}\left[X_{i}=k \cap X_{j}=k\right]=0$. However, $\operatorname{Pr}\left[X_{i}=k\right] \operatorname{Pr}\left[X_{j}=k\right]=$ $(1 / n)^{2 k}$. Thus, $X_{i}$ and $X_{j}$ are not independent.
(c) Let $M=\max \left(X_{1}, X_{2}, \ldots, X_{n}\right)$. Show $\operatorname{Pr}[M \geq 2 k / n] \leq n e^{-k /(3 n)}$.

Solution. The idea is to use Chernoff bounds to show that $\operatorname{Pr}\left[X_{i} \geq 2 k / n\right]$ is small and then use the union bound to bound the probability that any of the $X_{i}$ variables is greater than $2 k / n$. Recall that $X_{i}=\sum_{j=1}^{k} Y_{j i}$. We have $\operatorname{Pr}\left[X_{i} \geq(1+\delta) E\left[X_{i}\right]\right] \leq e^{-\delta^{2} E\left[X_{i}\right] / 3}$ by Chernoff. Thus, we get $\operatorname{Pr}\left[X_{i} \geq 2 k / n\right] \leq e^{-k /(3 n)}$, and by union bound $\operatorname{Pr}[M \geq 2 k / n] \leq$ $\sum_{i=1}^{n} \operatorname{Pr}\left[X_{i} \geq 2 k / n\right] \leq \sum_{i=1}^{n} e^{-k /(3 n)}=n e^{-k /(3 n)}$.

### 3.3 The Coupon Collector problem.

Suppose there are $n$ different kinds of coupons, and we want to collect at least one coupon from every kind. We start out with nothing, and at each step, we get a new random coupon, equally likely to be any of the $n$ kinds, and independent of the previous coupons. This is known as the coupon collector's problem.

- What is the expected time $T$ when we're done collecting?
- What is the variance of $T$ ?
- Use Chebyshev's inequality to bound the probability that $T$ deviates far from its expectation.

Solution. Let $T_{i}$ be the random variable equal to the first time we have $i$ different kinds of coupons. Then, we can break the total time to collect all kinds of coupons $T_{n}$ into the phases between getting a new kind of coupon:

$$
\begin{aligned}
\mathbb{E}\left[T_{n}\right] & =\mathbb{E}\left[T_{1}+\left(T_{2}-T_{1}\right)+\ldots+\left(T_{n}-T_{n-1}\right)\right] \\
& =\mathbb{E}\left[T_{1}\right]+\mathbb{E}\left[T_{2}-T_{1}\right]+\ldots+\mathbb{E}\left[T_{n}-T_{n-1}\right]
\end{aligned}
$$

Now let's think about the random variable $T_{k+1}-T_{k}$ : it is the time it takes us to get a $k+1$-th coupon given that we already have $k$ coupons. No matter what kinds of coupons we have already, the probability that we get a new coupon is $\frac{n-k}{n}$ in each step independently.
This is identical to the earlier problem where we had a Bernoulli random variable $X$ such that $\operatorname{Pr}[X=1]=p$, and we showed that the expected time until it becomes 1 for the first time is $\frac{1}{p}$. Thus, $\mathbb{E}\left[T_{k+1}-T_{k}\right]=\frac{n}{n-k}$, and

$$
\begin{aligned}
\mathbb{E}\left[T_{n}\right] & =1+\frac{n}{n-1}+\ldots+\frac{n}{1} \\
& =n\left(\frac{1}{n}+\frac{1}{n-1}+\ldots+\frac{1}{1}\right)=n H_{n}
\end{aligned}
$$

where $H_{n}$ is the $n$-th harmonic number. It is known that $H_{n}=\Theta(\log n)$ (which can be proved using an integral among other methods), hence $\mathbb{E}[T]=\Theta(n \log n)$.
For the variance, note that the random variables $T_{k+1}-T_{k}$ are independent. Indeed, if $k>l$, we have

$$
\operatorname{Pr}\left[T_{k+1}-T_{k}=t_{k}, T_{l+1}-T_{l}=t_{l}\right]=\operatorname{Pr}\left[T_{k+1}-T_{k}=t_{k} \mid T_{l+1}-T_{l}=t_{l}\right] \operatorname{Pr}\left[T_{l+1}-T_{l}=t_{l}\right]
$$

Now, note that conditioning on $T_{l+1}-T_{l}=t_{l}$ has no effect on the probability that $T_{k+1}-T_{k}=t_{k}$, since the future coupons we get are independent of the past. Hence the above is

$$
=\operatorname{Pr}\left[T_{k+1}-T_{k}=t_{k}\right] \operatorname{Pr}\left[T_{l+1}-T_{l}=t_{l}\right]
$$

which shows that the random variables are indeed independent. This means that

$$
\begin{aligned}
\operatorname{Var}\left[T_{n}\right] & =\operatorname{Var}\left[T_{1}+\left(T_{2}-T_{1}\right)+\ldots+\left(T_{n}-T_{n-1}\right)\right] \\
& =\operatorname{Var}\left[T_{1}\right]+\operatorname{Var}\left[T_{2}-T_{1}\right]+\ldots+\operatorname{Var}\left[T_{n}-T_{n-1}\right]
\end{aligned}
$$

Now we're faced with the general task of computing the variance of the random variable $T$ which is the first time that a Bernoulli random variable $X$ with $\operatorname{Pr}[X=1]=p$ becomes 1 . We have

$$
\operatorname{Pr}[T=t]=(1-p)^{t-1} p
$$

and as we saw earlier, $\mathbb{E}[T]=\frac{1}{p}$. It remains to compute

$$
\begin{aligned}
\mathbb{E}\left[T^{2}\right] & =\sum_{t=1}^{\infty} \operatorname{Pr}[T=t] t^{2} \\
& =\sum_{t=1}^{\infty}(1-p)^{t-1} p t^{2} \\
& =p \sum_{t=1}^{\infty}(1-p)^{t-1} t^{2}
\end{aligned}
$$

We could compute this sum by decomposing it into simpler sums in a clever way. But here's a useful (and more principled) trick for computing sums like this: consider the function $f(x)=\frac{1}{1-x}$ for $|x|<1$. Then we have the power series expansion

$$
\frac{1}{1-x}=1+x+x^{2}+\ldots=\sum_{n=0}^{\infty} x^{n}
$$

Differentiating both sides, we have

$$
\frac{1}{(1-x)^{2}}=1+2 x+3 x^{2}+\ldots=\sum_{t=0}^{\infty}(t+1) x^{t}
$$

and differentiating again,

$$
\frac{2}{(1-x)^{3}}=2+6 x+12 x^{2}+\ldots=\sum_{t=0}^{\infty}(t+1)(t+2) x^{t}
$$

Using this, we have

$$
\begin{aligned}
\sum_{t=1}^{\infty}(1-p)^{t-1} t^{2} & =\sum_{t=1}^{\infty}(1-p)^{t-1} t(t+1)-\sum_{t=1}^{\infty}(1-p)^{t-1} t \\
& =\sum_{t=0}^{\infty}(1-p)^{t}(t+1)(t+2)-\sum_{t=0}^{\infty}(1-p)^{t}(t+1) \\
& =\frac{2}{p^{3}}-\frac{1}{p^{2}}
\end{aligned}
$$

and so

$$
\begin{aligned}
\operatorname{Var}[T] & =\mathbb{E}\left[T^{2}\right]-\mathbb{E}[T]^{2} \\
& =\frac{2}{p^{2}}-\frac{1}{p}-\frac{1}{p^{2}}=\frac{1-p}{p^{2}}
\end{aligned}
$$

which implies that

$$
\operatorname{Var}\left[T_{n}\right]=\sum_{k=1}^{n} \frac{1-\frac{n-k}{n}}{\left(\frac{n-k}{n}\right)^{2}}=\sum_{k=1}^{n} \frac{n k}{(n-k)^{2}} \leq n^{2} \sum_{l=1}^{\infty} \frac{1}{l^{2}} \leq 2 n^{2}
$$

Thus, by Chebyshev,

$$
\operatorname{Pr}\left[\left|T_{n}-\mathbb{E}\left[T_{n}\right]\right| \geq c n\right] \leq \frac{2}{c^{2}}
$$


[^0]:    ${ }^{1}$ Formally, this means that $p(H)=p(T)$.

[^1]:    ${ }^{1}$ The Bernoulli distribution is defined by the random variable $X$ where $X=1$ with probability $p$ and $X=0$ with probability $1-p$.

